# 2Q: A Low Overhead High Performance Buffer Management Replacement Algorithm \*

**Paper Presentation** 

Author: Prayash Joshi

Department of Computer Science prayash@vt.edu





- Motivation
- 2Q
- **Experimental Results**
- **Tuning Parameters**
- Conclusion





#### Buffer Management and LRU Issues

- Buffer management is critical in minimizing disk I/O and improving throughput in database systems.
- The LRU (Least Recently Used) algorithm is widely used, but it suffers from high overhead and poor performance in certain workloads.
- LRU's performance deteriorates with sequential scans or random access to large files.
- Goal: Introduce an algorithm that maintains the simplicity of LRU but reduces the overhead and improves performance.





## 2Q Algorithm

- 20 is a self-tuning improvement over LRU and its variants (e.g., LRU/2).
- It maintains constant time overhead per page access (O(1)) while achieving better performance.
- Key mechanism: separating pages into two queues:<sup>1</sup>
  - **Al**: Holds pages on their first reference (FIFO queue).
  - **Am**: Holds "hot" pages (frequently accessed, LRU policy).
- Goal: Keep frequently accessed pages in the main buffer and discard cold pages quickly.





<sup>&</sup>lt;sup>1</sup>Iohnson and Shasha 1994.

#### How 2Q Works: Al and Am Queues

- When a page is first accessed, it is placed in the **Al** queue (FIFO).
- If it is accessed again while in Al, it is promoted to Am, the LRU queue.
- Pages in Am are considered "hot" and are retained for longer periods in the buffer.
- If not accessed again, the page is discarded from Al without promotion to Am.<sup>2</sup>





<sup>&</sup>lt;sup>2</sup>Iohnson and Shasha 1994.

#### Alin and Alout: Enhancements in 2Q

- 2Q introduces **Alin** and **Alout** for improved efficiency:
  - Alin: Holds recently accessed pages (pages with first access).
  - **Alout**: Holds references (tags) to evicted pages from Alin.
- Pages in Alout that are re-accessed are reintroduced to Am.
- This ensures sustained popularity tracking and prevents frequent cold pages from occupying buffer space.<sup>3</sup>





<sup>&</sup>lt;sup>3</sup>Johnson and Shasha 1994.

#### Performance of 2Q in Simulations

- Synthetic workloads and real-world traces were used for testing.
- Results:
  - 2Q consistently outperforms LRU and Gclock across all tests.
  - Matches or slightly exceeds LRU/2 in hit rate.
  - Shows a 5-10% improvement in hit rate over LRU.<sup>4</sup>





<sup>&</sup>lt;sup>4</sup>Johnson and Shasha 1994.

#### Real-World Tests: DB2 and UNIX Traces

- Real-world traces from DB2 commercial databases and UNIX systems.
- 2Q showed higher hit rates compared to LRU and Gclock.
- DB2 Results: Best performance when Kin (Alin size) was tuned to correlate with reference activity.
- 2Q required little tuning and remained efficient even with suboptimal parameters.<sup>5</sup>





<sup>&</sup>lt;sup>5</sup>Johnson and Shasha 1994.

#### Tuning Alout for Optimal Performance

- The size of **Alout** plays a critical role in responsiveness.
- Experimentally determined that setting Alout to around 50% of buffer size balances responsiveness and hit rate.
- Trade-off: A larger Alout improves responsiveness but may reduce long-term hit rate slightly.<sup>6</sup>





<sup>&</sup>lt;sup>6</sup>Iohnson and Shasha 1994.

#### Strengths of the 2Q Algorithm

- Constant time overhead per page access.
- Improved hit rates (5-10% over LRU).
- Adaptability to mixed workloads (sequential scans and random accesses).
- Efficient memory usage (tags stored in Alout).
- Minimal tuning required, performs well even in suboptimal conditions.





#### Limitations of the 2Q Algorithm

- Limited gains in extremely large buffers (performance plateaus).
- Inefficient in purely random access workloads.
- More complex than LRU due to management of Alin and Alout.
- Limited comparison to other modern algorithms like ARC.





## Conclusion: Key Takeaways

- 2Q provides an efficient alternative to LRU with constant time overhead and better performance in real-world workloads.
- It handles mixed access patterns efficiently, offering a balance between simplicity and performance.
- The self-tuning nature of 2Q makes it practical for high-performance database systems with minimal manual tuning required.





# Questions?



#### References L



lohnson, Theodore and Dennis Shasha (1994). "2Q: A Low Overhead High Performance Buffer Management Replacement Algorithm". In: Proceedings of the 20th International Conference on Very Large Data Bases. VLDB '94. San Francisco, CA, USA: Morgan Kaufmann Publishers Inc., 439-450. ISBN: 1558601538.

